

RECHERCHE

INSTITUT NATIONAL DES SCIENCES APPLIQUÉES DE LYON

Wireless Sensor Networks in a Nutshell

Séminaire Internet du Futur, ASPROM

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LYON



Agenda

- ✓ A short overview of the CITI lab
- ✓ Wireless Sensor Networks : Key applications & constraints
- ✓ Energy consumption and network *lifetime*
- ✓ Radio channel properties
- ✓ Key results
 - ✓ Resources sharing
 - ✓ Routing protocol and data gathering
- ✓ Conclusions & Open problems



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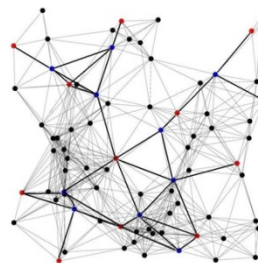
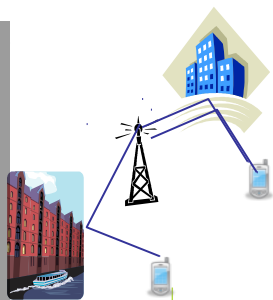
A short introduction to the CITI lab.



Details on: <http://citi.insa-lyon.fr/>

Research areas in the CITI lab.

- Software radio
- Green radio
- Cooperation & relaying
- RF architectures



- Multi-hop networks
- Self-networking
- Routing
- Green networking
- Real time

Networking

Radiocomms



Ambient networking

Security



- Secured protocols
- Trust mechanisms
- Code generation
- Privacy



Embedded systems



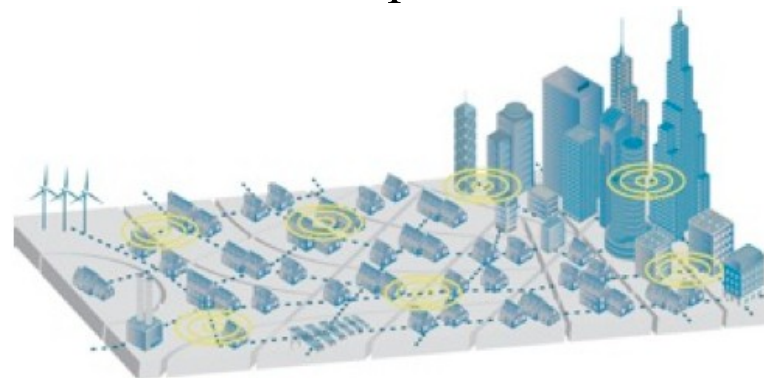
- SoC , NoC
- Software radio
- Simulation & testbeds
- sensors

Middleware



- Component based software platform
- OSGi
- Service oriented development

- ✓ Urbanet (leader: Dr. H. Rivano) focuses on
 - ✓ Context: Smart cities, digital societies
 - ✓ Focus on *capillary networks* (generally speaking: wireless sensor and actuator networks + wireless multi-hop mesh networks)
 - ✓ Goal: to provide networking optimization mechanisms and networking protocols to support ambient services

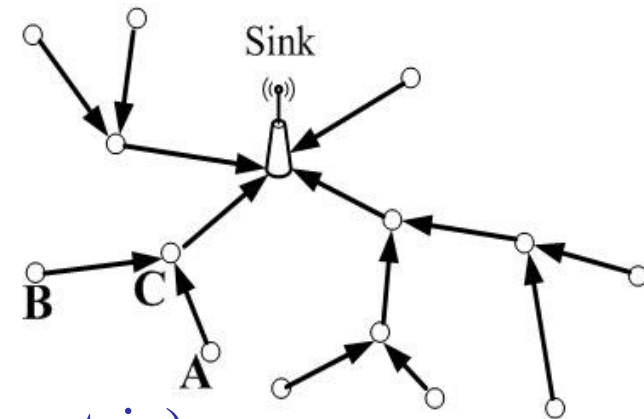




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Wireless Sensor Networks: Applications & Constraints



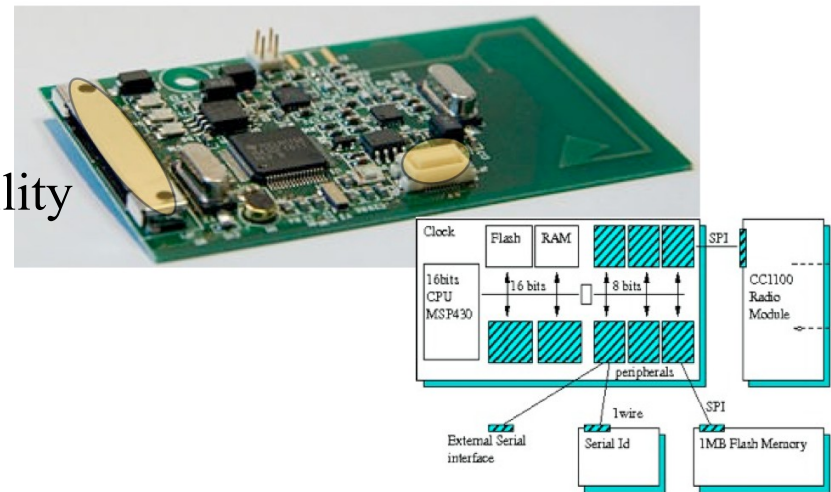
- ✓ Key entities for the Internet of Things
- ✓ Application-based networks (*aka* data-centric)
 - ✓ Physical measures using a physical sensor (water-metering, temperature control, etc.)
 - ✓ Coverage problem on a monitored area (intrusion detection, environment monitoring, wild animals tracking, etc.)
- ✓ Convergecast traffic to reach the sink node(s):
 - ✓ Alarms; periodical monitoring; request/response
 - ✓ Multi-hop paradigm from source to destination
 - ✓ Nodes to nodes traffic is limited

Wireless Sensor Networks: Applications & Constraints (cont'd)

- ✓ WSN networks topology properties
 - ✓ Random or regular (grid, line)
 - ✓ Network degree vary from 4/5 nodes (agricultural sensors) to thousand (urban networks for water-metering)
 - ✓ Network diameter varies from 3/4 hops to 10
 - ✓ Static nodes but the topology may be dynamic (to due sleeping mode, the volatility of the radio channel, etc.)

- ✓ Hardware properties

- ✓ Limited computation capability
- ✓ Low memory
- ✓ Embedded system
- ✓ Lifetime
- ✓ Low cost (low quality??)





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- ✓ Key issue: to maximize the network lifetime, defined as
 - ✓ Dead of the 1st wireless sensor node
 - ✓ or... Loss of connectivity between node(s) and the sink(s)
 - ✓ or... Coverage problem failed



→ Network lifetime = 10 years

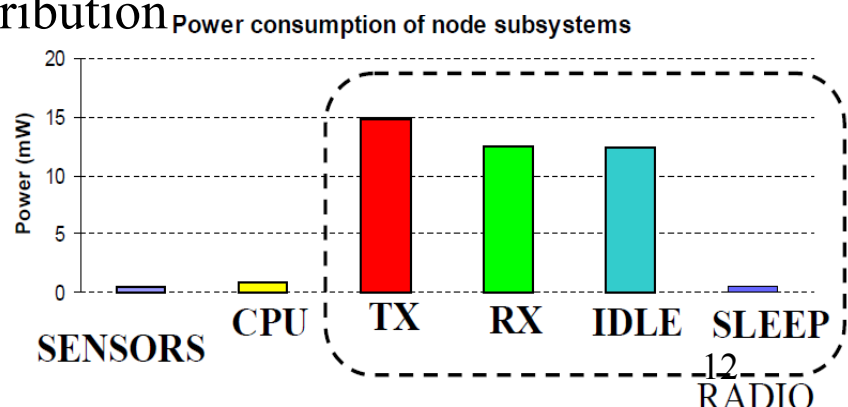
Energy issue (cont'd)

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→ Network lifetime = 10 years

- ✓ Focus on the radio transmission
 - In terms of energy consumption, to transmit 1 bit requires more than 1'000 CPU-cycles
 - Energy consumption distribution



- ✓ Network lifetime optimization:
 - ✓ Less for more!
(less transmission for more duration)

- ✓ All the opportunities we have:
 - ✓ Low energy consumption hardware system
 - ✓ Energy harvesting system
 - ✓ Energy-efficient radio interfaces
 - ✓ Sleeping mode for sensor nodes and efficient resource sharing
 - ✓ Energy-aware routing protocol (or, at least, energy-efficient routing protocol)
 - ✓ Data-aggregation

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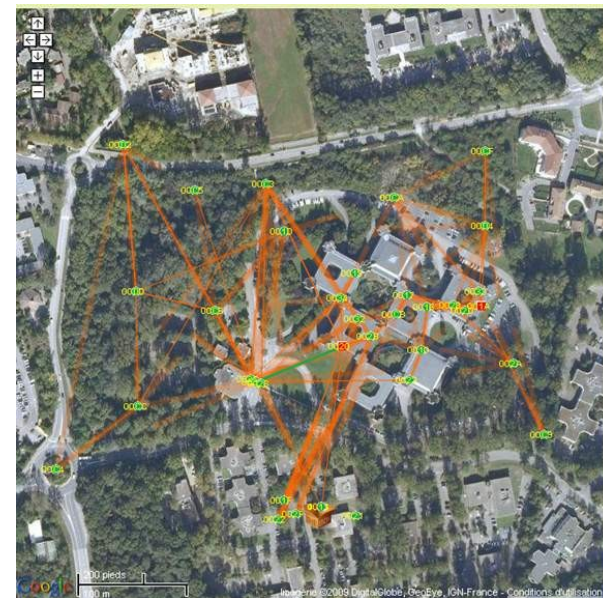
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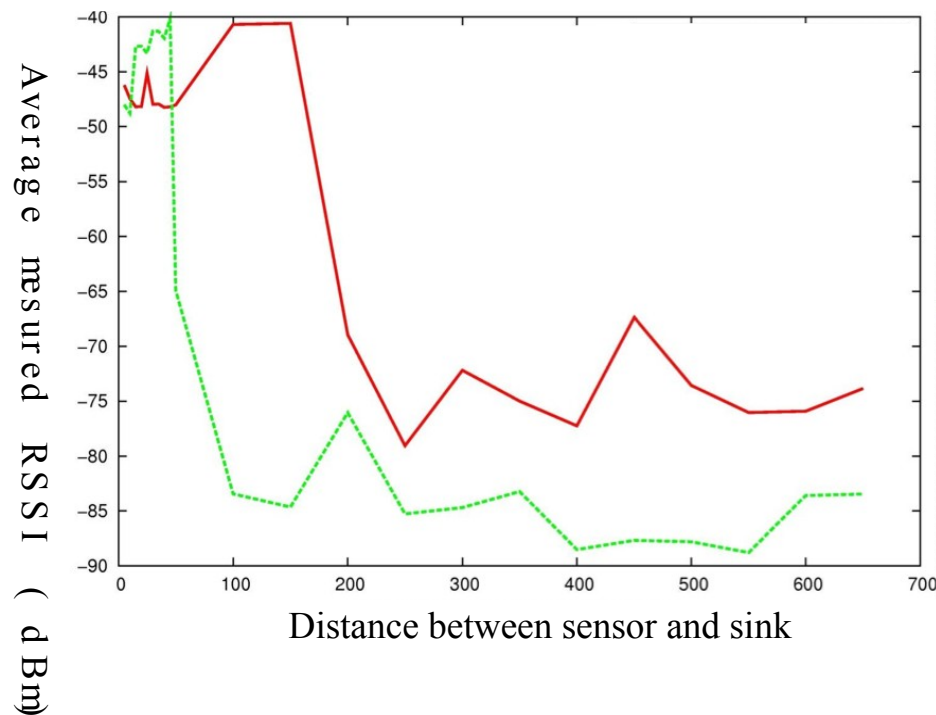
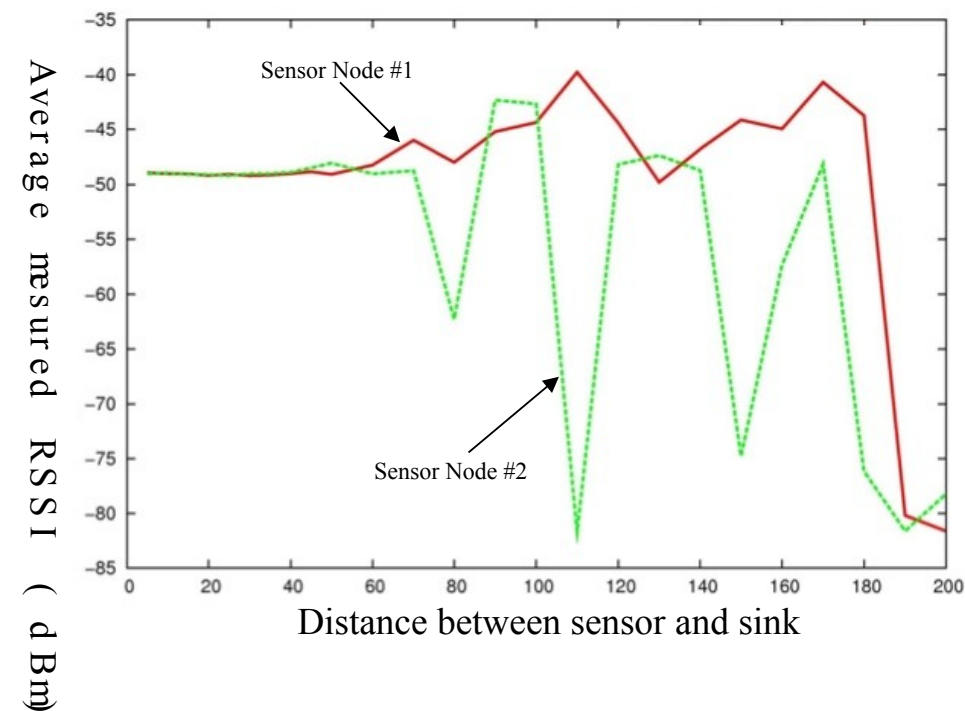
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- ✓ French project ANR ARESA (2006-2009):
 - ✓ More than 40 nodes (indoor/outdoor)
 - ✓ Trace with more than 400'000 packets
- ✓ Ph.D. of K. Heurtefeux (2009):
 - ✓ Apartment, CITI, soccer playground
- ✓ We investigate the RSSI behavior (Radio Strenght Indicator)
- ✓ What we have learned:
 - ✓ Results are material-dependent
 - ✓ Opportunistic radio links, asymmetric property
 - ✓ Radio channel is not stable in space and time
 - ✓ Other well-known phenomenon : fading, shadowing, interferences

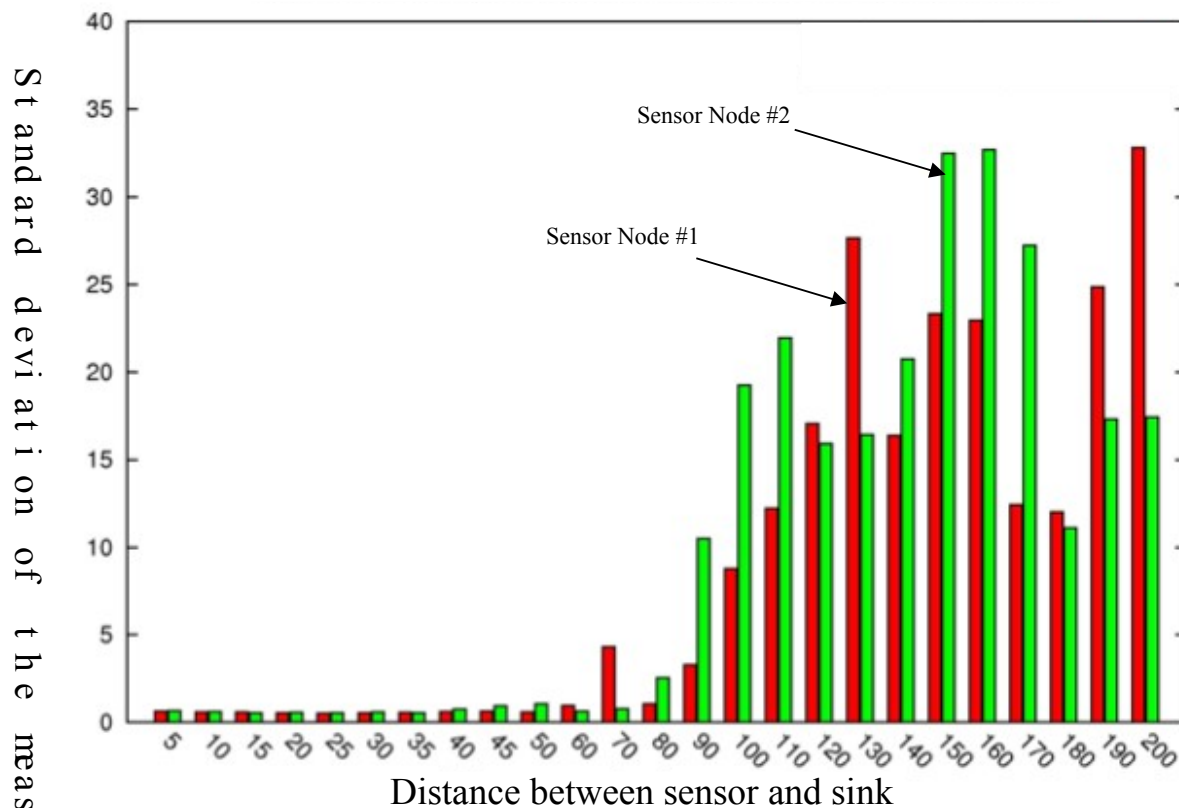


- ✓ Some RSSI exemples (apartment, CITI lab)
 - Hardware-dependent
 - Environment-dependent

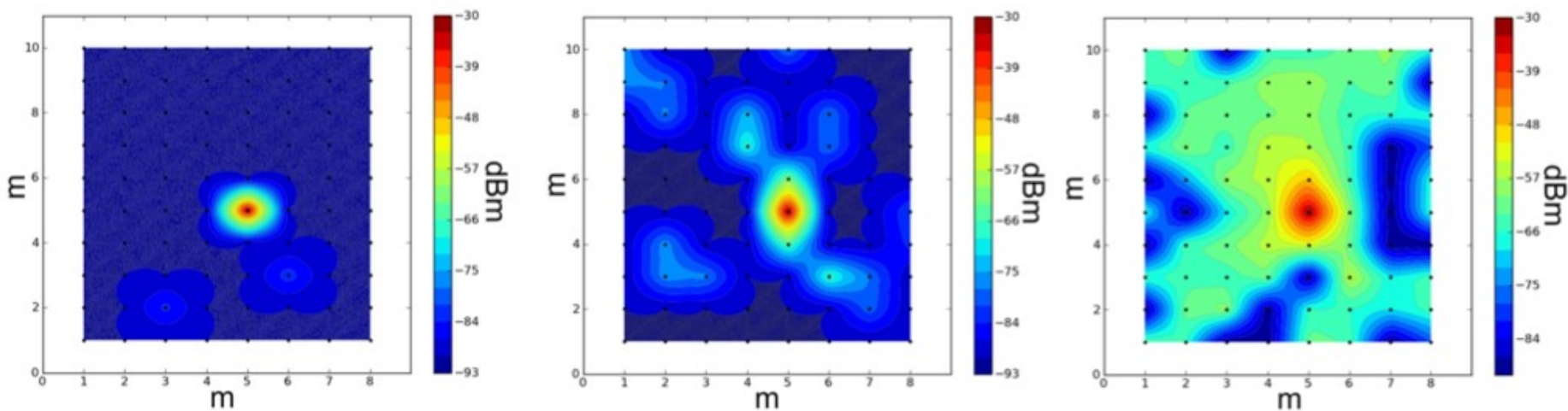
RSSI face to the distance



- ✓ RSSI variability (standard deviation)

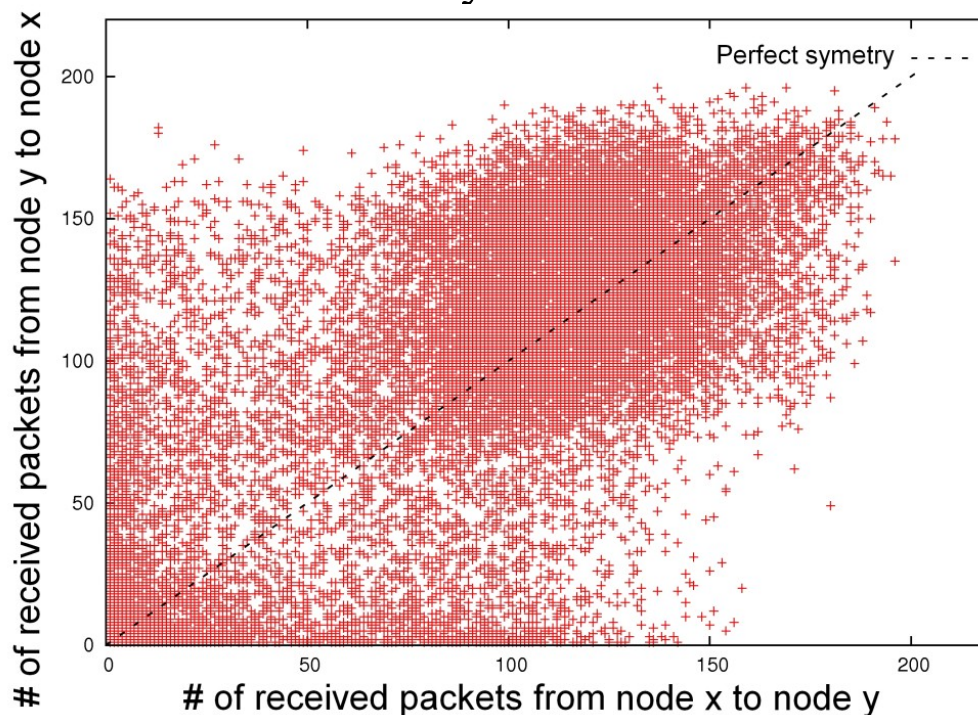


- ✓ Radio propagation is non-isotropic



Radio propagation on the SensLab testbed – Strasbourg site – (-30 dBm, -15 dBm, 0 dBm)

- ✓ Radio links are not always symmetric
 - Hardware-dependent, time-dependent, space-dependent
 - On the SensLab testbed (Grenoble site), more than 40% of radio links are non symmetric

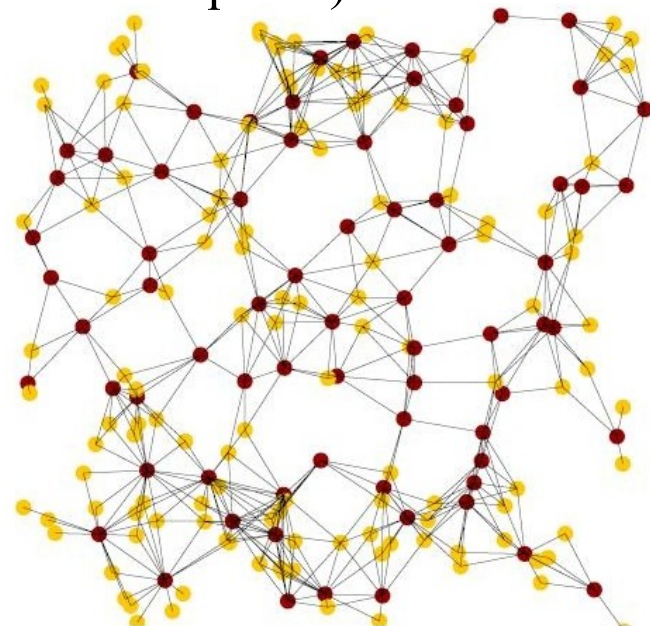




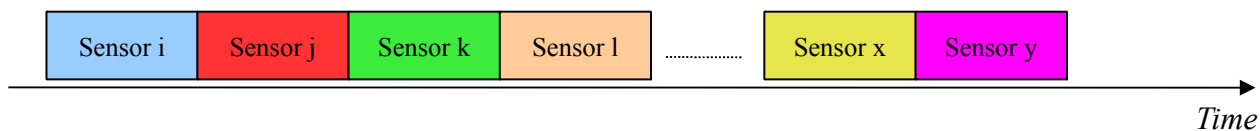
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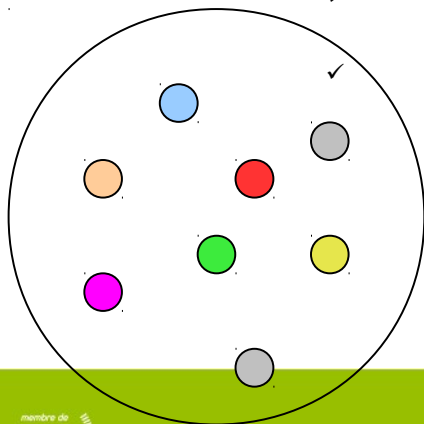
- ✓ MAC protocols (*Medium Access Control*)
- ✓ Goal: distributed and *fair* sharing of the radio channel using local information (1-hop neighborhood information), and with low collision probability
 - ✓ Deterministic Access (synchronisation is required)
 - ✓ Random Access (not necessarily using synchronisation)



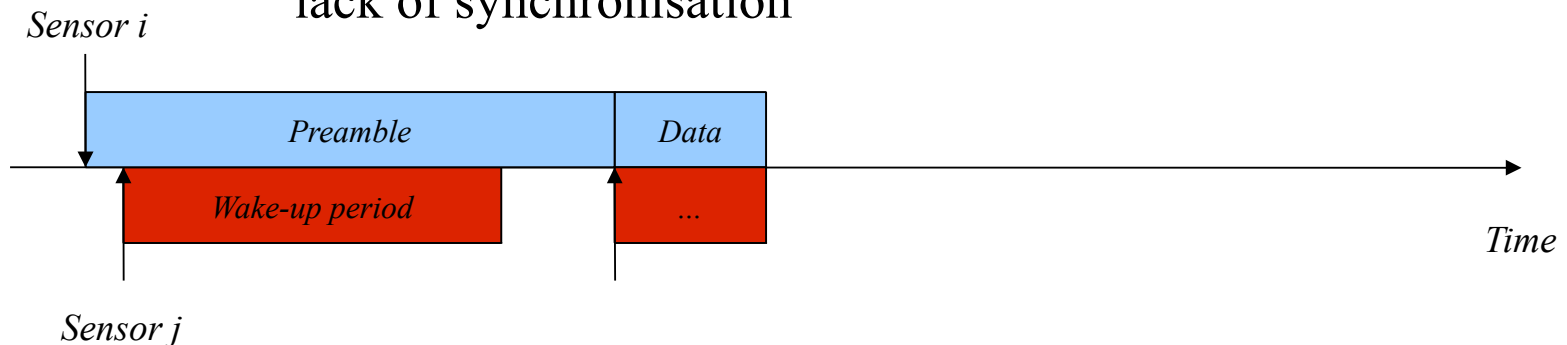
- ✓ **Deterministic access:**
 - ✓ Local scheduling is defined
 - ✓ Close to a TDMA approach (*Time Division Multiple Access*)
 - ✓ Each slot-time is allocated to a dedicated node



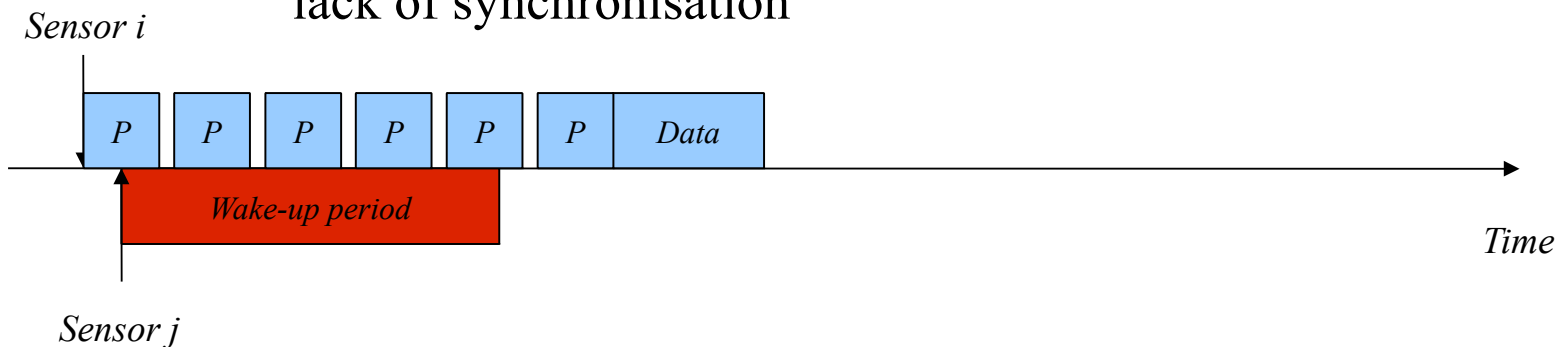
- ✓ Requires a fine synchronisation
- ✓ Non suitable for network dynamicity
- ✓ Not easy to cope with variable traffic intensity



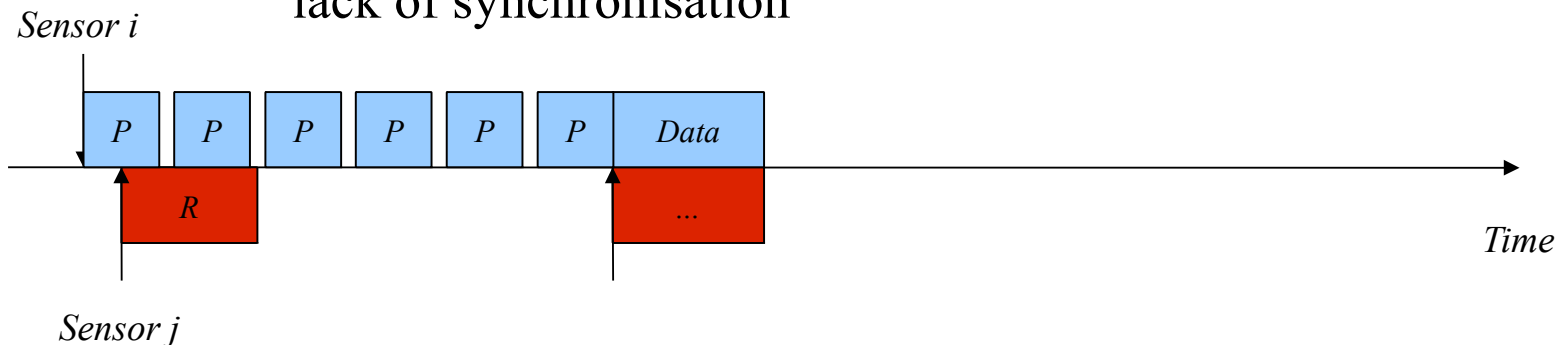
- ✓ Contention-based random access based
 - ✓ Based on a CSMA-like protocol but including sleeping mode for sensor nodes (duty-cycle mechanism)
 - ✓ 2 families : w/o Synchronisation & w/ Synchronisation
- ✓ Without synchronisation (BMAC, XMAC, ...)
 - ✓ Using preamble sampling strategy
 - ✓ Nodes wake up periodically but at different time due to the lack of synchronisation



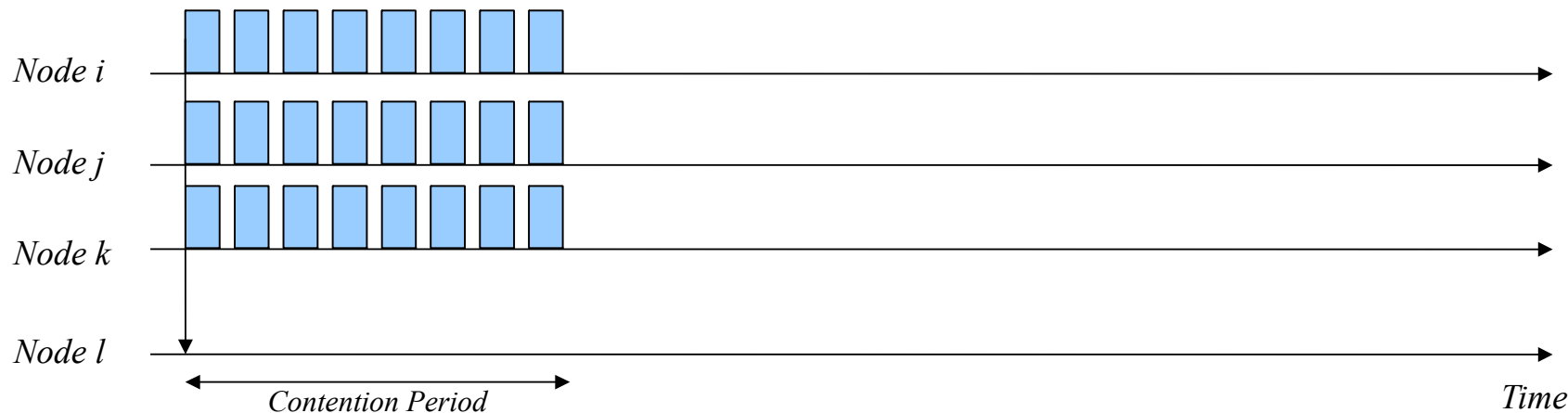
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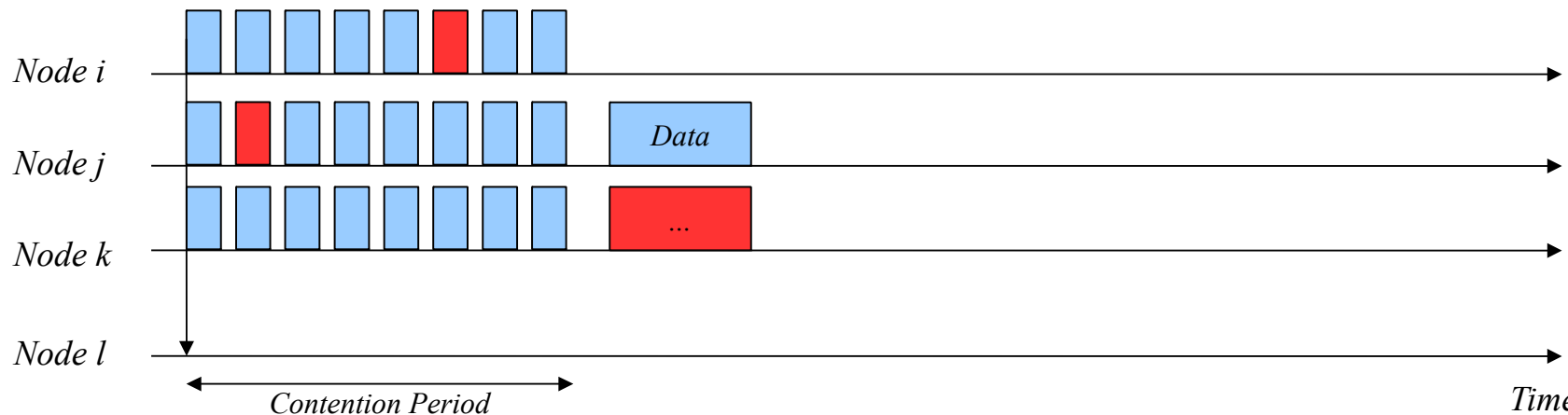
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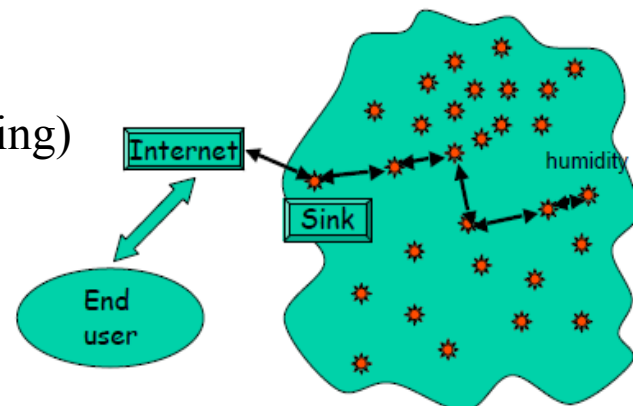
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 - ✓ Common clock
 - ✓ Periodical rendez-vous point



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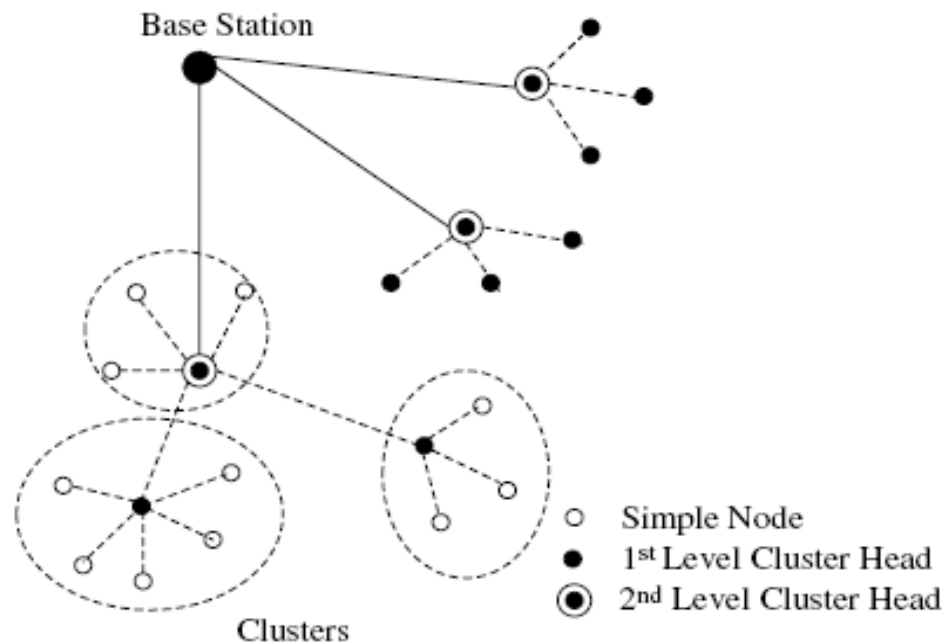


- ✓ Key idea: shortest path (in terms of either number of hops or euclidian distance or energy consumed)
- ✓ Some protocols come from mobile ad hoc networks (MANET)
 - ✓ But not really suitable because of too important overhead, huge *signalling* (periodical beacon and route management), energy wasting
- ✓ Dedicated protocols:
 - ✓ Hierarchical approaches
 - ✓ Location of Interests (content based routing)
 - ✓ Gradient-based routing protocols
 - ✓ Geographic (using GPS coordinates)
 - ✓ But also: multi-paths, QoS based, etc.



- ✓ Hierarchical approaches

- Using clusters, virtual backbone, cluster-tree, etc.



- ✓ Location of interests
 - Content-based routing protocols
 - Publish / subscribe policies

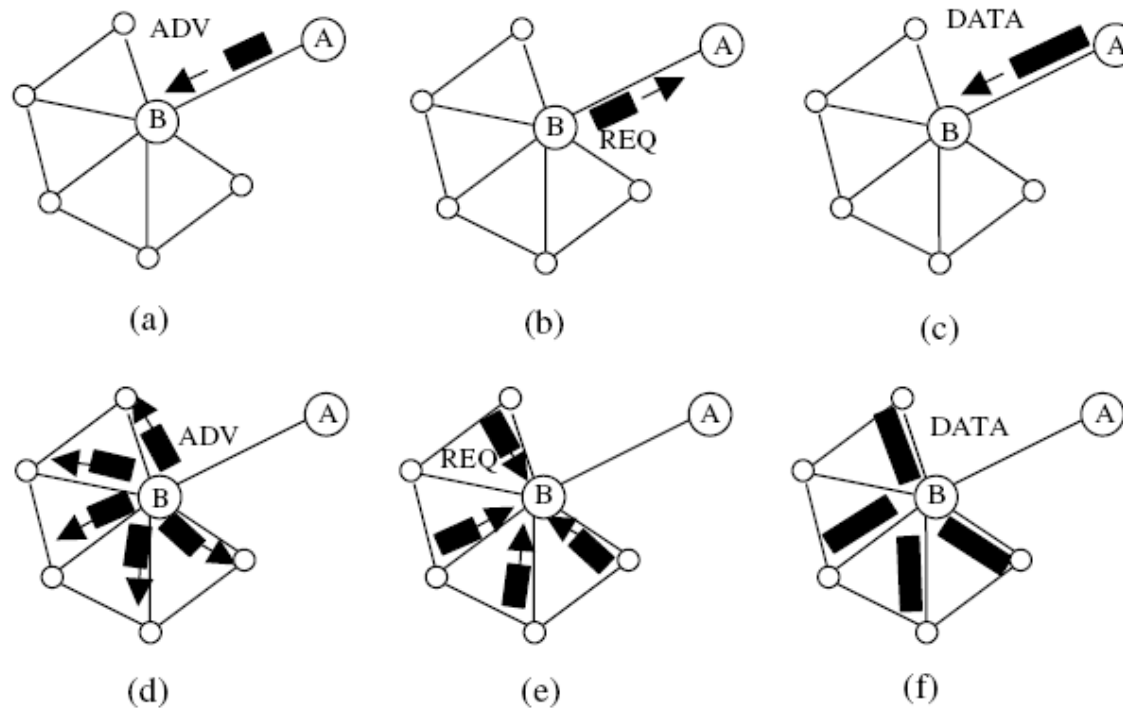
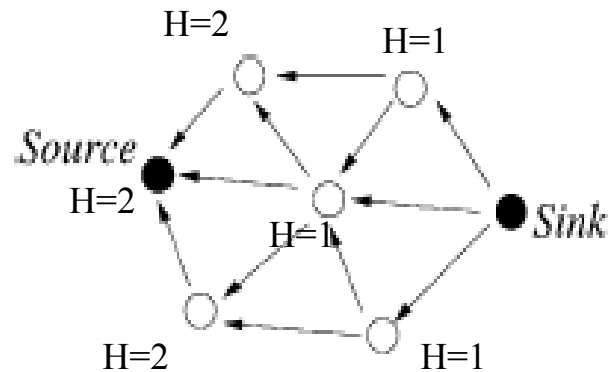


Fig. 3. SPIN protocol. Node A starts by advertising its data to node B (a). Node B responds by sending a request to node A (b). After receiving the requested data (c), node B then sends out advertisements to its neighbors (d), who in turn send requests back to B (e-f).

Routing protocol and data dissemination (cont'd)

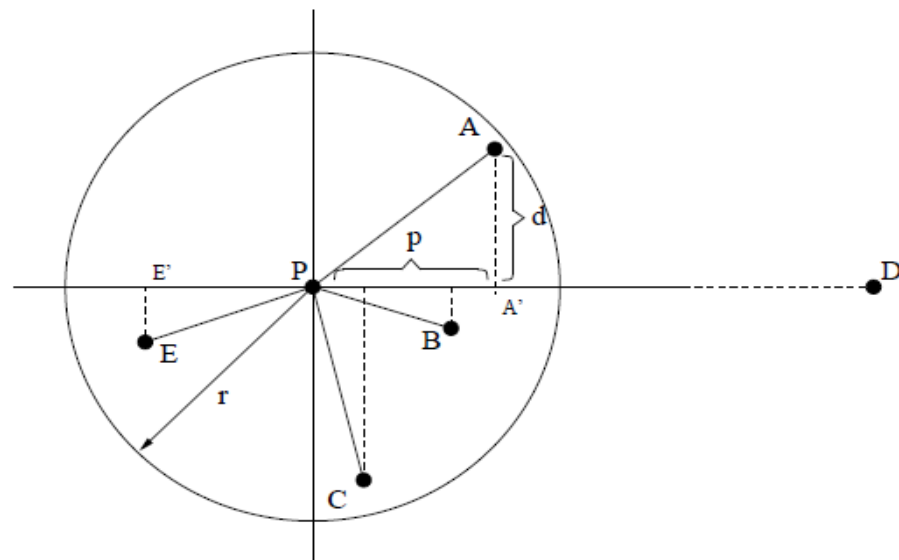
- ✓ Gradient routing protocol:
 - Flooding of an *init* packet, from the sink to the whole network
 - At each step, to increment the counter value



Routing protocol and data dissemination (cont'd)

✓ Geographic approaches

- ✓ Each node owns a unique Id. and a coordinate (x,y,z) "
- ✓ Absolute coordinates (GPS) or virtual coordinates
- ✓ Assume the sink location / sink coordinates
- ✓ Assume that a *well-known* function $f(x)$ exists such as: $f(Id.) \rightarrow (x, y)$
- ✓ The next forwarder is a neighbor which closer to the destination
 - Beacon-based (neighborhood is known a priori)
 - Beaconless (neighborhood is never known)

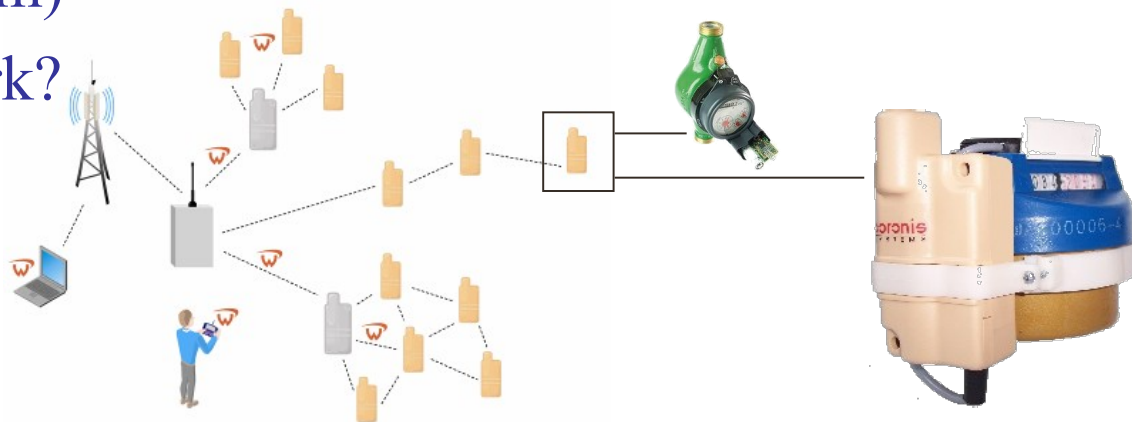




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- ✓ WSN are data-centric network
- ✓ Energy is the main challenge
 - ✓ Network lifetime optimization is a major concern
 - ✓ Cross-layer approaches (joint MAC/routage schemes)
- ✓ New issues: temporal constraints and QoS requirements
- ✓ To save energy: transmit less → data-aggregation
- ✓ Security (open system)
- ✓ IP-compliant network?
- ✓ From sensor nodes to dust...



Thank you for your attention,
Questions?

Contact: fabrice.valois@insa-lyon.fr
<http://fvalois.insa-lyon.fr/>